

What is a *bursty event*?

The burstiness of an event at any time t is defined by the *acceleration* of event mentionings for a defined time span of observation. *Burstiness* measures the increase in incoming rate of event mentionings, rather than the incoming rate itself. This measure quickly identifies important events that are emerging.

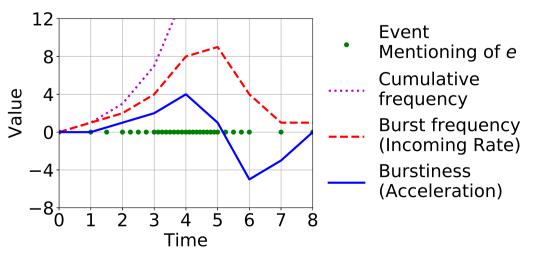
What is bursty event detection *throughout histories*?

Existing works focus on the detection of current trending events, but it is important to be able to go back in time and explore bursty events throughout the history. Our system can answer the following type of questions in realtime with space-efficient probabilistic data structure and algorithms.

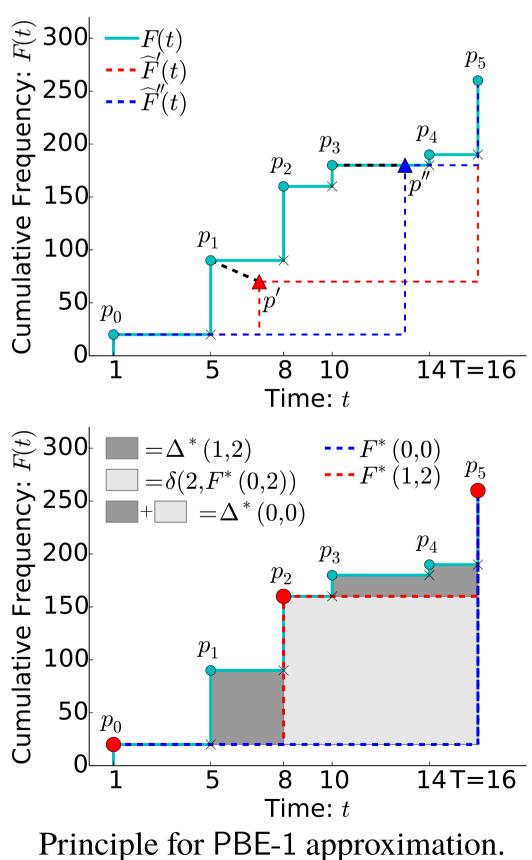
(i) What are the bursty events in the first week of October in 2016? (ii) Is "Anthem Protest" a bursty event in second week of September in 2017?

Notation	Meaning
e	an event or an event id
$M = \{(m_1, t_1), (m_2, t_2), \cdots \}$	stream of timestamped
	messages
$h: m_i \to [1, K]$	map a message to an even id
$S = \{(a_1, t_1), (a_2, t_2), \cdots \}$	event stream with
	(id, timestamp)
S_e	event stream of e
	(timestamp only)
$S[t_1, t_2]$	substream of S in time range
	$[t_1,t_2]$
τ	burst span
$F_e(t)$	cumulative frequency of e
	in $S[0, t]$
$f_e(t_1,t_2)$	frequency of e in $S[t_1, t_2]$
$bf_e(t) = f_e(t - \tau, t)$	burst frequency (incoming
	rate) of e at t
$b_e(t) = bf_e(t) - bf_e(t - \tau)$	burstiness (acceleration)
	of e at t

Preliminaries



An example of burst where $\tau = 1$.



PBE-1 Buffered



Key Contributions

1. The paper first presents a formal definition of burstiness b(t) for any historical time t. 2. PBE: The probabilistic data structure estimates burstiness b(t) in a single event stream. **3.** Two PBE's, PBE-1: Approximation with buffering, PBE-2: Approximation without buffering. 4. CM-PBE: Estimate burstiness $b_e(t)$ of any event e at any time t over a multi-event stream. 5. Bursty Event Detection Query: Find all the bursty events for a given a time range [start, end]. 6. Extensive empirical investigation using Olympic Rio 2016 and US Election 2016 tweets. 7. Space efficient and query efficient probabilistic data structure with theoretical guarantees.

PBE-2 Non-Buffered CM-PBE & Results $\int 000 \, F(t)$ $\frac{F(t)}{Error}$ $PBE_{1,h1(e)}$:, 250 eduency $h_1(e)$ $PBE_{2,h2}(e)$ $\mathbf{1}\gamma$ $h_2(e)$ (e, t) h₃(e) <u>፲</u> 150 h₄(e) Cumulative 05 00 $\overline{\text{PBE}}_{3,h3}(e)$ $\overline{\text{PBE}}_{4,h4(e)}$ $F(4) - \gamma$ Pictorial representation CM-PBE 14T = 1678910 45 Time: t (1) 300 $P_{F(t)}$ ← CM-PBE-1 **—** CM-PBE-2 100Error 150 Ц Cumulative 00 00 80 100 80 60 0178910 14T = 164 5 Space Usage (MB) Time: t

Principle for PBE-2 approximation.

Absolute Error vs Space OlympicRio

